

CSE 392: Matrix and Tensor Algorithms for Data

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Lecture 5: Matrix factorizations II - eigenvalue decomposition, PCA

Outline

1 Eigenvalue problems

2 PCA

3 Eigenfaces

Eigenvalue problems

Given a square matrix $\mathbf{A} \in \mathbb{C}^{n \times n}$, the *eigenvalue problem*:

$$\mathbf{A}\mathbf{u} = \lambda\mathbf{u}.$$

λ is an *eigenvalue* and \mathbf{u} is an *eigenvector* of \mathbf{A} .

Types of problems:

- Find the largest or the smallest eigenvalues.
- Compute all eigenvalues in region of \mathbb{C} .
- Compute dominant eigenvalues and eigenvectors.

Applications: Structural engineering, stability analysis, electronic structure calculations, dimensionality reduction, spectral clustering and graphs, pagerank and many more.

Eigenvalues and properties

A complex scalar λ is called an *eigenvalue* of a square matrix $\mathbf{A} \in \mathbb{C}^{n \times n}$ if there exists a nonzero vector $\mathbf{u} \in \mathbb{C}^n$ such that

$$\mathbf{A}\mathbf{u} = \lambda\mathbf{u}.$$

The vector \mathbf{u} is called an *eigenvector* of \mathbf{A} associated with λ .

- λ is an eigenvalue iff the columns of $\mathbf{A} - \lambda\mathbf{I}$ are linearly dependent.
- That is, $\det(\mathbf{A} - \lambda\mathbf{I}) = 0$.

Eigenvalues and properties II

- The set of all eigenvalues of \mathbf{A} , denoted $\Lambda(\mathbf{A})$, is the *spectrum* of \mathbf{A} .
- An eigenvalue is a root of the *characteristic polynomial*:

$$p_{\mathbf{A}}(\lambda) = \det(\mathbf{A} - \lambda\mathbf{I})$$

- So there are n eigenvalues (counted with their multiplicities).
- The multiplicity of these eigenvalues as roots of $p_{\mathbf{A}}$ are called *algebraic multiplicities*.
- The *geometric multiplicity* of an eigenvalue λ_i is the number of linearly independent eigenvectors associated with λ_i .
- Geometric multiplicity is \leq algebraic multiplicity.

Eigenvalues and properties III

- **Diagonalization:** Two matrices \mathbf{A} , \mathbf{B} are *similar* if there exists a nonsingular matrix \mathbf{X} such that: $\mathbf{A} = \mathbf{X}\mathbf{B}\mathbf{X}^{-1}$.
- $\mathbf{A}\mathbf{u} = \lambda\mathbf{u} \Leftrightarrow \mathbf{B}(\mathbf{X}^{-1}\mathbf{u}) = \lambda(\mathbf{X}^{-1}\mathbf{u})$
eigenvalues remain the same, eigenvectors transformed.
- \mathbf{A} is diagonalizable if it is similar to a diagonal matrix.
- *Transformations that preserve eigenvectors:*
 - ▶ Shift : $\mathbf{B} = (\mathbf{A} - \eta\mathbf{I})$
 - ▶ Polynomial : $\mathbf{B} = p(\mathbf{A})$
 - ▶ Inverse: $\mathbf{B} = \mathbf{A}^{-1}$
 - ▶ Shift and inverse: $\mathbf{B} = (\mathbf{A} - \eta\mathbf{I})^{-1}$

Symmetric eigenvalue problem

- For every square symmetric matrix $\mathbf{A} \in \mathbb{R}^{n \times n}$, we can compute eigendecomposition:

$$\mathbf{A} = \mathbf{U}\mathbf{\Lambda}\mathbf{U}^\top,$$

where \mathbf{U} is an orthogonal matrix with eigenvectors \mathbf{u}_i as columns, and $\mathbf{\Lambda}$ is diagonal matrix with eigenvalues λ_i on the diagonal.

- \mathbf{U} forms an orthonormal basis of eigenvectors of \mathbf{A} .
- Eigenvalues of \mathbf{A} are real.
- When \mathbf{A} is real, \mathbf{U} is also real.

The min-max theorem (Courant-Fischer)

Label eigenvalues decreasingly: $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$.

The eigenvalues of a Hermitian matrix \mathbf{A} are characterized by the relation

$$\lambda_k = \max_{\mathbf{S}, \dim(\mathbf{S})=k} \min_{\mathbf{x} \in \mathbf{S}, \mathbf{x} \neq 0} \frac{\langle \mathbf{A}\mathbf{x}, \mathbf{x} \rangle}{\langle \mathbf{x}, \mathbf{x} \rangle}$$

or

$$\lambda_k = \min_{\mathbf{S}, \dim(\mathbf{S})=n-k+1} \max_{\mathbf{x} \in \mathbf{S}, \mathbf{x} \neq 0} \frac{\langle \mathbf{A}\mathbf{x}, \mathbf{x} \rangle}{\langle \mathbf{x}, \mathbf{x} \rangle}$$

- $\frac{\langle \mathbf{A}\mathbf{x}, \mathbf{x} \rangle}{\langle \mathbf{x}, \mathbf{x} \rangle}$ is called the Rayleigh–Ritz quotient of \mathbf{A} .
- $\lambda_1 = \max_{\mathbf{x} \neq 0} \frac{\langle \mathbf{A}\mathbf{x}, \mathbf{x} \rangle}{\langle \mathbf{x}, \mathbf{x} \rangle}$ and $\lambda_n = \min_{\mathbf{x} \neq 0} \frac{\langle \mathbf{A}\mathbf{x}, \mathbf{x} \rangle}{\langle \mathbf{x}, \mathbf{x} \rangle}$.

Question: Use min-max theorem to show that $\sigma_1 = \|\mathbf{A}\|_2$.

Interlacing Theorem

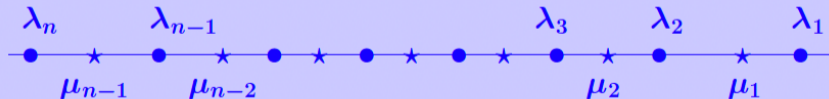
Suppose $\mathbf{A} \in \mathbb{R}^{n \times n}$ is symmetric. Let $\mathbf{B} \in \mathbb{R}^{m \times m}$ with $m < n$ be a principal submatrix (obtained by deleting both i -th row and i -th column for some values of i).

Suppose \mathbf{A} has eigenvalues $\lambda_1 \geq \dots \geq \lambda_n$, and \mathbf{B} has eigenvalues $\mu_1 \geq \dots \geq \mu_m$. Then

$$\lambda_k \geq \mu_k \geq \lambda_{n+k-m} \text{ for } k = 1, \dots, m$$

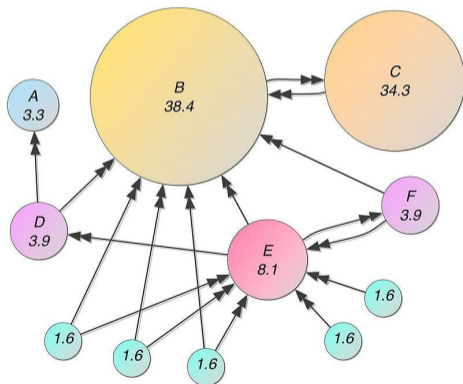
and if $m = n - 1$,

$$\lambda_1 \geq \mu_1 \geq \lambda_2 \geq \mu_2 \geq \dots \geq \mu_{n-1} \geq \lambda_n$$



PageRank

- **PageRank** is the first Google algorithm developed to evaluate the quality and importance of web pages.
- Webgraph - created by all World Wide Web pages as nodes and hyperlinks as edges.
- Likelihood that a person randomly clicking on links will arrive at any particular page.



- PageRank value of a page is given as:

$$PR(p_i) = \frac{1-d}{N} + d \sum_{p_j \in M(p_i)} \frac{PR(p_j)}{L(p_j)},$$

p_1, p_2, \dots, p_N are the pages, $M(p_i)$ = set of pages that link to p_i , $L(p_j)$ = number of outbound links on page p_j , N = total number of pages, and d = damping factor.

- The values are the entries of the dominant right eigenvector of the modified adjacency matrix rescaled so that each column adds up to one.

$$\mathbf{r} = \begin{bmatrix} PR(p_1) \\ PR(p_2) \\ \vdots \\ PR(p_N) \end{bmatrix}$$

- \mathbf{r} is the solution of the equation

$$\mathbf{r} = \begin{bmatrix} (1-d)/N \\ (1-d)/N \\ \vdots \\ (1-d)/N \end{bmatrix} + d \begin{bmatrix} \ell(p_1, p_1) & \ell(p_1, p_2) & \cdots & \ell(p_1, p_N) \\ \ell(p_2, p_1) & \ddots & & \vdots \\ \vdots & & \ell(p_i, p_j) & \\ \ell(p_N, p_1) & \cdots & & \ell(p_N, p_N) \end{bmatrix} \mathbf{r}$$

the adjacency function $\ell(p_i, p_j)$ is the ratio between number of links outbound from page j to page i to the total number of outbound links of page j .

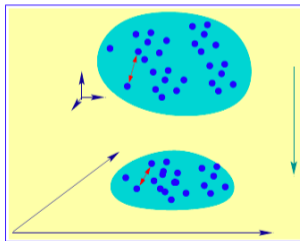
- $$\sum_{i=1}^N \ell(p_i, p_j) = 1,$$

The matrix is a stochastic matrix. Closely related to the problem of finding the stationary points of Markov processes. It is also a variant of the eigenvector centrality measure used commonly in network analysis.

Dimensionality Reduction

Dimensionality Reduction

- **Dimensionality Reduction** (DR) techniques pervasive to many data applications.
- Reduce computational cost; but also more often :
 - ▶ reduce noise and redundancy in data, and
 - ▶ discover patterns.
- Given $\mathbf{x} \in \mathbb{R}^d$, and $k \ll d$, find the mapping $\Phi : \mathbf{x} \in \mathbb{R}^d \longrightarrow \mathbf{y} \in \mathbb{R}^k$.



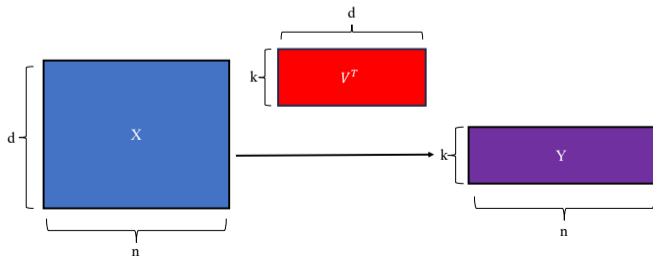
Projection-based Dimensionality Reduction

- Given dataset $\mathbf{X} = [\mathbf{x}_1, \dots, \mathbf{x}_n]$, and dimension k , find the reduced set \mathbf{Y} .
- *Projection method*: Explicit mapping to the lower dimension

$$\mathbf{y} = \mathbf{V}^\top \mathbf{x}$$

with $\mathbf{V} \in \mathbb{R}^{d \times k}$.

- Projection-based Dimensionality Reduction : $\mathbf{Y} = \mathbf{V}^\top \mathbf{X}$. Find the best such mapping (optimization) given that the \mathbf{y}_i 's must satisfy certain constraints.



Principal Component Analysis

- **Principal Component Analysis (PCA)** : find (orthogonal) \mathbf{V} so that projected data $\mathbf{Y} = \mathbf{V}^\top \mathbf{X}$ has maximum variance.
- Maximize over all orthogonal $d \times k$ matrices \mathbf{V} :

$$\sum_i \|\mathbf{y}_i - \frac{1}{n} \sum_j \mathbf{y}_j\|_2^2 = \dots = \text{Tr}[\mathbf{V}^\top \bar{\mathbf{X}} \bar{\mathbf{X}}^\top \mathbf{V}],$$

where $\bar{\mathbf{X}} = [\bar{\mathbf{x}}_1, \dots, \bar{\mathbf{x}}_n]$ with $\bar{\mathbf{x}}_i = \mathbf{x}_i - \boldsymbol{\mu}$, and $\boldsymbol{\mu}$ = mean.

- *Solution*: \mathbf{V} = dominant k eigenvectors of the covariance matrix. Top k singular vectors of $\bar{\mathbf{X}}$.

Exercises

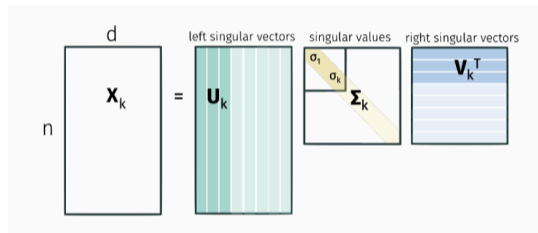
- Show that $\bar{\mathbf{X}} = \mathbf{X}(\mathbf{I} - \frac{1}{n}\mathbf{e}\mathbf{e}^\top)$ (here \mathbf{e} = vector of all ones). What does the projector $(\mathbf{I} - \frac{1}{n}\mathbf{e}\mathbf{e}^\top)$ do?
- Show that solution \mathbf{V} also minimizes reconstruction error:

$$\sum_i \|\bar{\mathbf{x}}_i - \mathbf{V}\mathbf{V}^\top \bar{\mathbf{x}}_i\|^2 = \sum_i \|\bar{\mathbf{x}}_i - \mathbf{V}\bar{\mathbf{y}}_i\|^2$$

- It also maximizes $\sum_{i,j} \|\mathbf{y}_i - \mathbf{y}_j\|^2$

Low rank approximation

- Given a data matrix $\mathbf{X} \in \mathbb{R}^{n \times d}$ and integer k , find a rank- k approximation of \mathbf{X} .
- $\mathbf{X}_k = \mathbf{U}_k \mathbf{\Sigma}_k \mathbf{V}_k^\top = \mathbf{U}_k \mathbf{U}_k^\top \mathbf{X} = \mathbf{X} \mathbf{V}_k \mathbf{V}_k^\top$.



$$\mathbf{U}_k = \arg \min_{\mathbf{U} \in \mathbb{R}^{n \times k}} \|\mathbf{X} - \mathbf{U} \mathbf{U}^\top \mathbf{X}\|_F^2 = \arg \max_{\mathbf{U} \in \mathbb{R}^{n \times k}} \|\mathbf{U} \mathbf{U}^\top \mathbf{X}\|_F^2.$$

$$\|\mathbf{X} - \mathbf{X}_k\|_F^2 = \sum_{i=k+1}^n \sigma_i^2.$$

Eigenfaces