

CSE 392: Matrix and Tensor Algorithms for Data

Instructor: Shashanka Ubaru

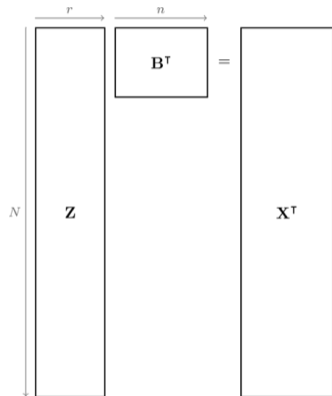
University of Texas, Austin
Spring 2024

Lecture 18: Randomized CP - II

① CP-ARLS-Mix

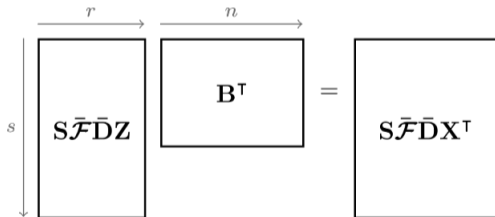
② CP-ARLS-Lev

$$\min_{\mathbf{B}} \|\mathbf{Z}\mathbf{B}^\top - \mathbf{X}^\top\|_F^2$$



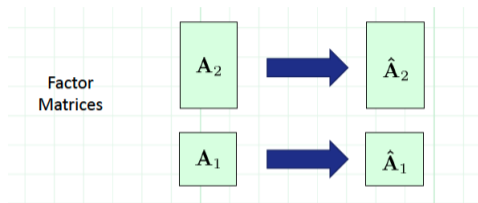
$$\min_{\mathbf{B}} \|\mathbf{S}\bar{\mathcal{F}}\bar{\mathbf{D}}\mathbf{Z}\mathbf{B}^\top - \mathbf{S}\bar{\mathcal{F}}\bar{\mathbf{D}}\mathbf{X}^\top\|_F^2$$

- \mathbf{S} is $s \times N$ sampling matrix
- $\bar{\mathcal{F}} = \mathcal{F}_d \otimes \cdots \otimes \mathcal{F}_{k+1} \otimes \mathcal{F}_{k-1} \otimes \cdots \otimes \mathcal{F}_1$.
- $\bar{\mathbf{D}} = \mathbf{D}_d \otimes \cdots \otimes \mathbf{D}_{k+1} \otimes \mathbf{D}_{k-1} \otimes \cdots \otimes \mathbf{D}_1$.



$$\mathbf{Z} = \mathbf{A}_d \odot \cdots \odot \mathbf{A}_{k+1} \odot \mathbf{A}_{k-1} \odot \cdots \odot \mathbf{A}_1$$

Mixing KRP Efficiently Using Kronecker FJLT



$$\begin{aligned} S\bar{F}\bar{D}Z &= S(\mathcal{F}_2 \otimes \mathcal{F}_1)(D_2 \otimes D_1)(A_2 \odot A_1) \\ &= S((\mathcal{F}_2 D_2) \otimes (\mathcal{F}_1 D_1))(A_2 \odot A_1) \\ &= S((\mathcal{F}_2 D_2 A_2) \odot (\mathcal{F}_1 D_1 A_1)) \\ &= S(\hat{A}_2 \odot \hat{A}_1) \end{aligned}$$

Pre-Mixing Tensor

Need to compute sketched right hand side...

$$S\bar{F}\bar{D}X^\top = S(\mathcal{F}_2 \otimes \mathcal{F}_1)(D_2 \otimes D_1)X_{(3)}^\top$$

Pre-mixed tensor

$$\tilde{\mathcal{X}} = \mathcal{X} \times_1 \mathcal{F}_1 D_1 \times_2 \mathcal{F}_2 D_2 \times_3 \mathcal{F}_3 D_3$$

$$\tilde{X}_{(3)}^\top = (\mathcal{F}_2 D_2 \otimes \mathcal{F}_1 D_1) X_{(3)}^\top (\mathcal{F}_3 D_3)^\top$$

Sample before unmixing

$$S\bar{F}\bar{D}X^\top = (S\tilde{X}_{(3)}^\top) D_3 \mathcal{F}_3^*$$

$$\min_B \|S\bar{F}\bar{D}ZB^\top - S\bar{F}\bar{D}X^\top\|_F^2$$

The diagram shows the matrix equation $S\bar{F}\bar{D}ZB^\top = S\bar{F}\bar{D}X^\top$. The matrix $S\bar{F}\bar{D}Z$ is a vertical rectangle with height s and width r . The matrix B^\top is a square with width n . The matrix $S\bar{F}\bar{D}X^\top$ is a vertical rectangle with height s and width n . Arrows indicate the dimensions: r for the width of $S\bar{F}\bar{D}Z$, n for the width of B^\top , and s for the height of $S\bar{F}\bar{D}Z$ and $S\bar{F}\bar{D}X^\top$.

CP-ARLS-Mix Algorithm

Inputs: Tensor $\mathcal{X} \in \mathbb{R}^{n_1 \times n_2 \times \dots \times n_d}$, desired rank $r \in \mathbb{N}$, number of samples $s \in \mathbb{N}$.

- 1 Initialize $\mathbf{A}_k \in \mathbb{R}^{n_k \times r}$ for all $k \in [d]$
- 2 Draw random diagonal \mathbf{D}_k for all $k \in [d]$
- 3 Compute $\tilde{\mathbf{A}}_k = \mathcal{F}_k \mathbf{D}_k \mathbf{A}_k$ for all $k \in [d]$
- 4 Compute $\tilde{\mathcal{X}} = \mathcal{X} \times_1 \mathcal{F}_1 \mathbf{D}_1 \times_2 \mathcal{F}_2 \mathbf{D}_2 \cdots \times_d \mathcal{F}_d \mathbf{D}_d$
- 5 $\Omega \leftarrow$ sampled indices for function value estimation
- 6 **repeat**
- 7 **for** $k = 1, \dots, d$ **do**
- 8 $\mathbf{S} \leftarrow$ random rows of \mathbf{I} scaled by $1/\sqrt{s}$.
- 9 $\hat{\mathbf{Z}} \leftarrow$ SKRP($\mathbf{S}, \tilde{\mathbf{A}}_1, \dots, \tilde{\mathbf{A}}_{k-1}, \tilde{\mathbf{A}}_{k+1}, \dots, \tilde{\mathbf{A}}_d$)
- 10 $\hat{\mathbf{X}} \leftarrow \mathcal{F}_k^* \mathbf{D}_k \left(\text{STU}(\mathbf{S}, \tilde{\mathcal{X}}, k) \right)$
- 11 $\mathbf{A}_k \leftarrow \arg \min_B \|\hat{\mathbf{Z}} \mathbf{B}^\top - \hat{\mathbf{X}}^\top\|_F^2$
- 12 $\tilde{\mathbf{A}}_k \leftarrow \mathcal{F}_k \mathbf{D}_k \mathbf{A}_k$
- 13 **end**
- 14 **until** SFV($\Omega, \mathcal{X}, \mathbf{A}_1, \mathbf{A}_2, \dots, \mathbf{A}_d$) ceases to decrease

Is the KFJLT adequate? YES

$$\min_B \|S\bar{\mathcal{F}}\bar{\mathbf{D}}\mathbf{Z}\mathbf{B}^\top - S\bar{\mathcal{F}}\bar{\mathbf{D}}\mathbf{X}^\top\|_F^2$$

- \mathbf{S} is $s \times N$ sampling matrix
- $\bar{\mathcal{F}} = \mathcal{F}_d \otimes \cdots \otimes \mathcal{F}_{k+1} \otimes \mathcal{F}_{k-1} \otimes \cdots \otimes \mathcal{F}_1$.
- $\bar{\mathbf{D}} = \mathbf{D}_d \otimes \cdots \otimes \mathbf{D}_{k+1} \otimes \mathbf{D}_{k-1} \otimes \cdots \otimes \mathbf{D}_1$.

- R. Jin, T. G. Kolda, and R. Ward. *Faster Johnson Lindenstrauss Transforms via Kronecker Products*, Information and Inference, 2020
- O. A. Malik, and S. Becker. *Guarantees for the Kronecker Fast Johnson Lindenstrauss Transform Using a Coherence and Sampling Argument*, Linear Algebra and its Applications, 2020

Recall: JL Lemma

JL Lemma

Let $\Phi \in \mathbb{R}^{m \times N}$ have independent entries $s_{ij} \sim \frac{1}{\sqrt{m}}\mathcal{N}(0, 1)$. If $m = O\left(\frac{\log(p)}{\epsilon^2}\right)$, then for any set of n data points $\mathbf{x}_1, \dots, \mathbf{x}_p \in \mathbb{R}^N$, with high probability:

$$(1 - \epsilon)\|\mathbf{x}_i - \mathbf{x}_j\|_2 \leq \|\Phi\mathbf{x}_i - \Phi\mathbf{x}_j\|_2 \leq (1 + \epsilon)\|\mathbf{x}_i - \mathbf{x}_j\|_2$$

Recall: JL Lemma

JL Lemma

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The Fast JL Lemma

Let $\Phi = \mathbf{SHD} \in \mathbb{R}^{m \times N}$ be a subsampled randomized Hadamard transform with $m = O\left(\frac{\log(N) \log(p)}{\epsilon^2}\right)$. Then for any set of p data points $\mathbf{x}_1, \dots, \mathbf{x}_p \in \mathbb{R}^N$, with high probability,

$$\|\Phi \mathbf{x}_i\|_2 = (1 \pm \epsilon) \|\mathbf{x}_i\|_2.$$

KFJL Result

Let $\Phi = \mathbf{S}(\mathcal{F}_d \mathbf{D}_d \otimes \cdots \otimes \mathcal{F}_1 \mathbf{D}_1) \in \mathbb{R}^{m \times N}$ be a KFJLT with $m = O\left(\frac{\log(N) \log^{2d-1}(p)}{\epsilon^2}\right)$.

Then for any set of p data points $\mathbf{x}_1, \dots, \mathbf{x}_p \in \mathbb{R}^N$, with high probability,

$$\|\Phi \mathbf{x}_i\|_2 = (1 \pm \epsilon) \|\mathbf{x}_i\|_2.$$

- R. Jin, T. G. Kolda, and R. Ward. *Faster Johnson Lindenstrauss Transforms via Kronecker Products*, Information and Inference, 2020

KFJLT-Sketch and solve

Given a matrix $\mathbf{A} \in \mathbb{R}^{N \times r}$ and a fixed vector $\mathbf{b} \in \mathbb{R}^N$, let $\mathbf{x}^* = \min_{\mathbf{x} \in \mathbb{R}^d} \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2$. Let $\Phi = \mathcal{S}(\mathcal{F}_d \mathbf{D}_d \otimes \cdots \otimes \mathcal{F}_1 \mathbf{D}_1) \in \mathbb{R}^{m \times N}$ be a KFJLT with

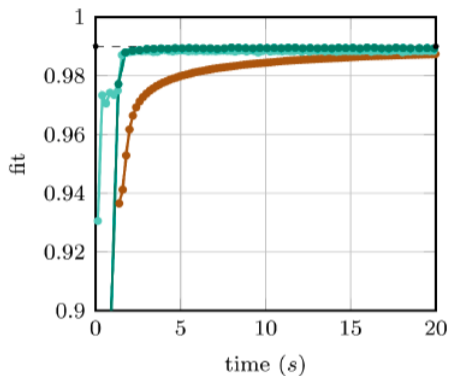
$$m = O\left(\frac{r^{2d} \log(N) \log^{2d-1}(r)}{\epsilon}\right),$$

and if $\tilde{\mathbf{x}} = \min_{\mathbf{x} \in \mathbb{R}^d} \|\Phi(\mathbf{A}\mathbf{x} - \mathbf{b})\|_2$, then, with high probability,

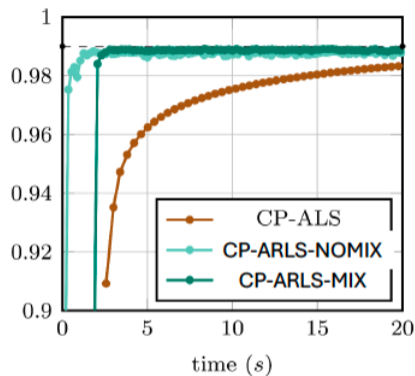
$$\|\mathbf{A}\tilde{\mathbf{x}} - \mathbf{b}\|_2 \leq (1 + \epsilon) \|\mathbf{A}\mathbf{x}^* - \mathbf{b}\|_2.$$

- R. Jin, T. G. Kolda, and R. Ward. *Faster Johnson Lindenstrauss Transforms via Kronecker Products*, Information and Inference, 2020

CP-ARLS faster than CP-ALS



300 x 300 x 300 Random Rank-5 Tensor
with 1% Noise



80 x 80 x 80 x 80 Random Rank-5 Tensor
with 1% Noise

Leverage scores and incoherence

Leverage scores

Given $\mathbf{A} \in \mathbb{R}^{N \times r}$, and an orthonormal basis \mathbf{U} for $\text{span}(\mathbf{A})$, for $i \in [n]$, the i th *leverage score*

$$\ell_i(\mathbf{A}) = \sup_{\mathbf{x}} \frac{(\mathbf{A}_{i*} \mathbf{x})^2}{\|\mathbf{A} \mathbf{x}\|^2} = \|\mathbf{U}_{i*}\|^2.$$

Leverage scores and incoherence

Leverage scores

Given $\mathbf{A} \in \mathbb{R}^{N \times r}$, and an orthonormal basis \mathbf{U} for $\text{span}(\mathbf{A})$, for $i \in [n]$, the i th *leverage score*

$$\ell_i(\mathbf{A}) = \sup_{\mathbf{x}} \frac{(\mathbf{A}_{i*} \mathbf{x})^2}{\|\mathbf{A} \mathbf{x}\|^2} = \|\mathbf{U}_{i*}\|^2.$$

Coherence

The coherence of $\mathbf{A} \in \mathbb{R}^{N \times r}$, denoted by $\mu(\mathbf{A})$ is the *maximum leverage score*, i.e.,

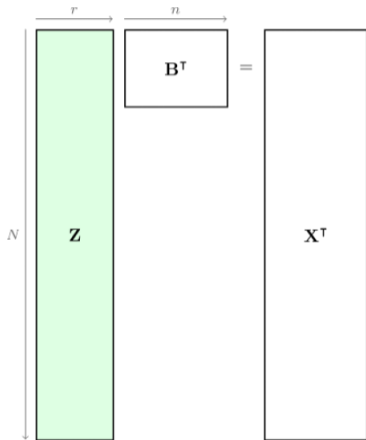
$$\mu(\mathbf{A}) = \max_{i \in [N]} \ell_i(\mathbf{A}).$$

We have $\frac{r}{N} \leq \mu(\mathbf{A}) \leq 1$.

We say \mathbf{A} is **incoherent** if $\mu(\mathbf{A}) \approx \frac{r}{N}$.

Is the KRP incoherent?

$$\min_{\mathbf{B}} \|\mathbf{Z}\mathbf{B}^\top - \mathbf{X}^\top\|_F^2$$



Khatri-Rao Product:

$$\mathbf{Z} = \mathbf{A}_d \odot \cdots \odot \mathbf{A}_{k+1} \odot \mathbf{A}_{k-1} \odot \cdots \odot \mathbf{A}_1$$

- Lemma 1: $\mu(\mathbf{A} \otimes \mathbf{B}) = \mu(\mathbf{A})\mu(\mathbf{B})$
- Lemma 2: $\mu(\mathbf{A} \odot \mathbf{B}) \leq \mu(\mathbf{A})\mu(\mathbf{B})$

KRP is incoherent if the factor matrices are!

Recall: Leverage score sampling

Sampling for LS

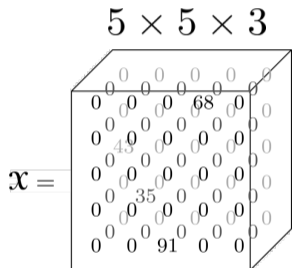
Given a matrix $\mathbf{A} \in \mathbb{R}^{N \times r}$ and a fixed vector $\mathbf{b} \in \mathbb{R}^N$, let $\mathbf{x}^* = \min_{\mathbf{x} \in \mathbb{R}^d} \|\mathbf{Ax} - \mathbf{b}\|_2$. Let $\mathbf{S} \in \mathbb{R}^{m \times N}$ be a sampling matrix with probabilities $p_i = \ell_i/r$, and $\mathbf{S}_{i*} = \mathbf{e}_j / \sqrt{mp_j}$ with $\Pr(j = i) = p_i$. If $m = O(r \log(r/\delta)/\epsilon)$ and $\tilde{\mathbf{x}} = \min_{\mathbf{x} \in \mathbb{R}^d} \|\mathbf{S}(\mathbf{Ax} - \mathbf{b})\|_2$, then, with high probability,

$$\|\mathbf{A}\tilde{\mathbf{x}} - \mathbf{b}\|_2 \leq (1 + \epsilon)\|\mathbf{Ax}^* - \mathbf{b}\|_2.$$

We also saw a procedure for approximately estimating the leverage scores.

Sparse Tensors

We can store a sparse tensor in size proportion to its number of nonzeros nnz



$$N = \prod_{k=1}^3 n_k = 75$$

$q = 4$ nonzeros

$$\mathbf{C} = \begin{bmatrix} 1 & 4 & 1 \\ 3 & 1 & 3 \\ 4 & 2 & 2 \\ 5 & 3 & 1 \end{bmatrix} \in \mathbb{N}^{q \times d} \quad \text{and} \quad \mathbf{v} = \begin{bmatrix} 68 \\ 43 \\ 35 \\ 91 \end{bmatrix} \in \mathbb{R}^q,$$

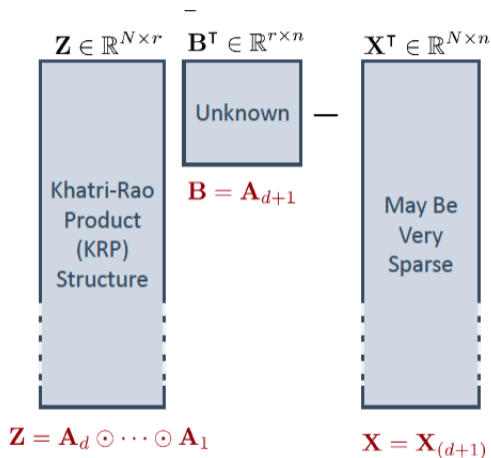
CP for sparse tensors

$$N \gg r, n$$

Linking back to mode-($d+1$) least squares subproblem

$$N = \prod_{k=1}^d n_k$$

$$n = n_{d+1}$$



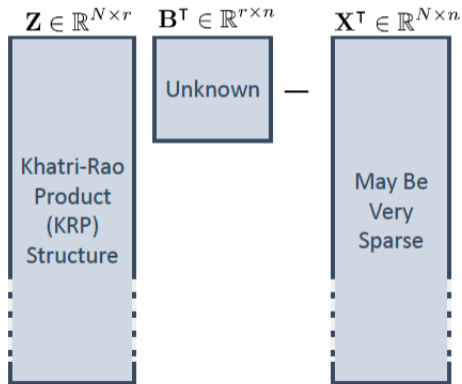
- KRP costs $O(Nr)$ to form
- System costs $O(Nnr^2)$ to solve
- KRP structure
 - Cost reduced to $O(Nnr)$
- KRP structure + data sparse
 - Cost reduced to $O(r \text{nnz}(\mathbf{X}))$

For Reddit (mode 3)
 $N = 1.2\text{B}$
 $\text{nnz}(\mathbf{X}) = 4.7\text{B}$

$$\min_B \|\mathbf{Z}\mathbf{B}^T - \mathbf{X}^T\|_F^2$$

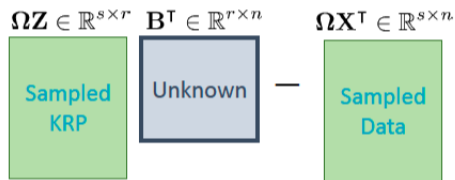
CP by sampling

$$\min_B \| \mathbf{ZB}^\top - \mathbf{X}^\top \|_F^2$$



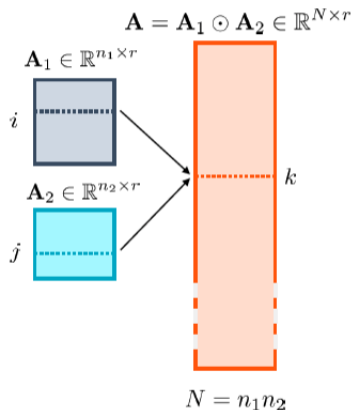
$$N \gg r, n$$

$$\min_B \| \Omega \mathbf{ZB}^\top - \Omega \mathbf{X}^\top \|_F^2$$



Complexity reduced from $O(Nnr)$ to $O(snr^2)$

Bounding Leverage Scores



Upper Bound on Leverage Score

Lemma (Cheng et al., NIPS 2016;

Battaglini et al., SIMAX 2018):

$$l_k(\mathbf{A}) \leq l_i(\mathbf{A}_1) l_j(\mathbf{A}_2)$$

Too expensive to calculate $\mathcal{O}(Nr^2)$

Cheap to calculate leverage scores for each submatrix $\mathcal{O}((n_1 + n_2)r^2)$

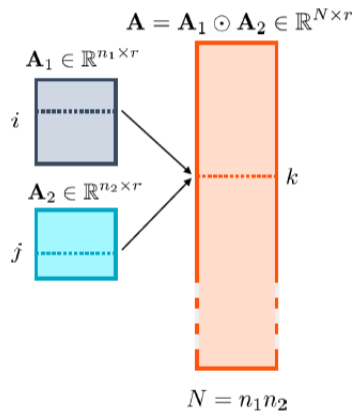
1-1 Correspondence between k and (i, j)

$$k \in [N] \Leftrightarrow (i, j) \in [n_1] \otimes [n_2]$$

Probability of Sampling row k in \mathbf{A} :

$$p_k = \frac{l_i(\mathbf{A}_1) l_j(\mathbf{A}_2)}{r^2}$$

Sampling Piecemeal



Upper Bound on Leverage Score

Lemma (Cheng et al., NIPS 2016;
Battaglini et al., SIMAX 2018):

$$\ell_k(\mathbf{A}) \leq \ell_i(\mathbf{A}_1) \ell_j(\mathbf{A}_2)$$

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Probability of Sampling row k in \mathbf{A} :

$$p_k = \frac{\ell_i(\mathbf{A}_1) \ell_j(\mathbf{A}_2)}{r^2}$$

Choose $i \sim p_i = \ell_i(\mathbf{A}_1)/r$

Choose $j \sim p_j = \ell_j(\mathbf{A}_2)/r$

$$k = i + (j - 1)n_1$$

Accuracy of Sketched ALS for 3-way Tensors

Original system with N rows

$$\mathbf{X}_{\text{opt}} = \arg \min_{\mathbf{X}} \|\mathbf{A}\mathbf{X} - \mathbf{B}\|_F^2$$

Sampled system with s rows

$$\tilde{\mathbf{X}}_{\text{opt}} = \arg \min_{\mathbf{X}} \|\tilde{\mathbf{A}}\mathbf{X} - \tilde{\mathbf{B}}\|_F^2$$

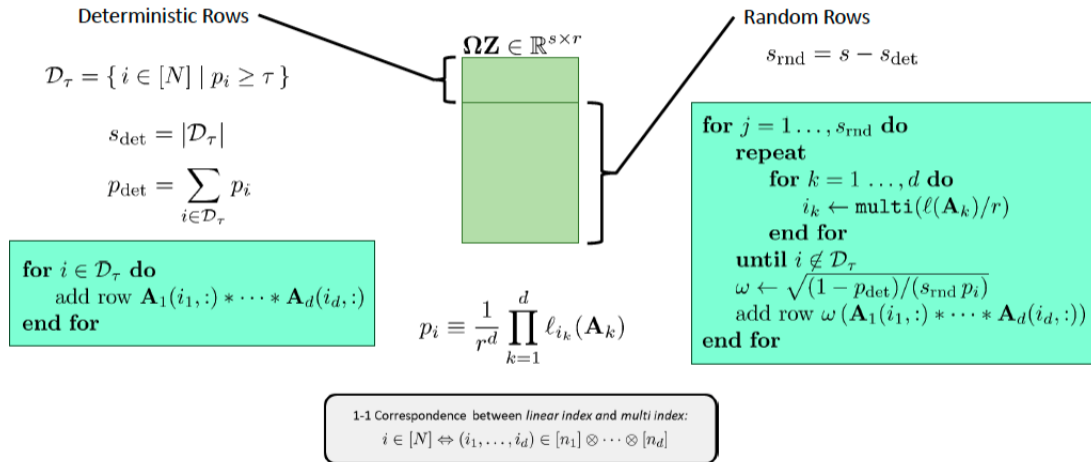
$$\text{Prob} \left(\|\mathbf{A}\tilde{\mathbf{X}}_{\text{opt}} - \mathbf{B}\|_F^2 \leq (1 + \varepsilon) \|\mathbf{A}\mathbf{X}_{\text{opt}} - \mathbf{B}\|_F^2 \right) > (1 - \delta)$$

$$\text{if } s = \frac{r}{\beta} \max \left\{ C \log \left(\frac{r}{\delta} \right), \frac{1}{\delta \varepsilon} \right\} \quad \text{where } \beta = \frac{1}{r} \leq \min_i \frac{p_i r}{\ell_i(\mathbf{A})} \in (0, 1]$$

$$\Rightarrow s = r^2 \max \left\{ C \log \left(\frac{r}{\delta} \right), \frac{1}{\delta \varepsilon} \right\}$$

Larsen & Kolda, SIAM J. Matrix Analysis & Applications (2022)

Hybrid Deterministic and Randomly Sampled Rows



Find All High Probability Rows without Computing All Probabilities

- Recall

$$p_i \equiv \frac{1}{r^d} \prod_{k=1}^d \ell_{i_k}(\mathbf{A}_k)$$

- For given tolerance $\tau > 1/N$, define the set of deterministic rows to include

$$\mathcal{D}_\tau = \{ i \in [N] \mid p_i \geq \tau \}$$

- Compute *without* computing all p_i values
- A few high leverage scores means all the others are necessarily low!
- Use bounding procedure to eliminate most options
- Compute products of at most a top few leverage scores in each mode

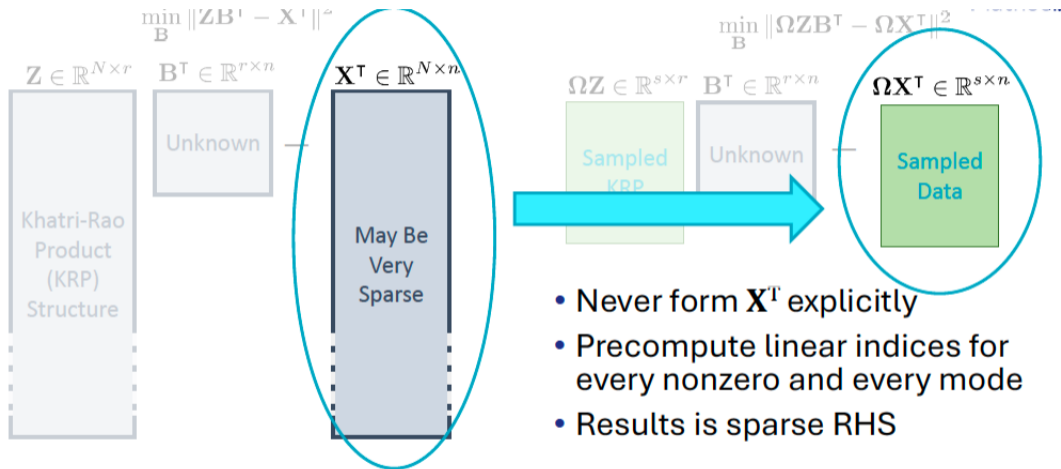
Sorted Leverages Scores (Descending)



1-1 Correspondence between linear index and multi index:

$$i \in [N] \Leftrightarrow (i_1, \dots, i_d) \in [n_1] \otimes \dots \otimes [n_d]$$

Efficiently Extract RHS from (Sparse) tensor



Similar in spirit to ideas for dense tensors in Battaglino et al., SIMAX 2018

Algorithm 3 CP via Alternating Randomized Least Squares with Leverage Scores

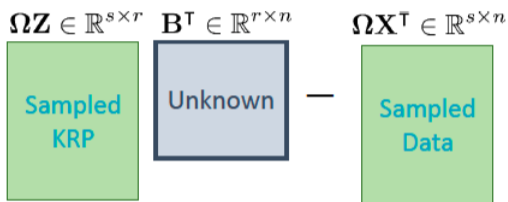
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1: function CP-ARLS-LEV( $\mathcal{X}$ ,  $r$ ,  $s$ ,  $\tau$ ,  $\eta$ ,  $\pi$ ,  $\text{tol}$ ,  $\{\mathbf{A}_k\}$ )
2:   for  $k = 1, \dots, d+1$  do
3:      $\mathbf{p}_k \leftarrow \ell(\mathbf{A}_k)/r$  ▷ Compute scaled leverage scores for initial guess
4:   end for
5:   repeat
6:     for  $\ell = 1, \dots, \eta$  do ▷ Group outer iterations into epochs
7:       for  $k = 1, \dots, d+1$  do ▷ Cycle through tensor modes
8:          $(\text{idx}, \text{wgt}, \bar{s}) \leftarrow \text{SKRPLEV}(\mathbf{p}_1, \dots, \mathbf{p}_{k-1}, \mathbf{p}_{k+1}, \dots, \mathbf{p}_{d+1}, s, \tau)$  ▷  $\bar{s} \leq s$ 
9:          $\tilde{\mathbf{Z}} \leftarrow \text{KRPSAMP}(\mathbf{A}_1, \dots, \mathbf{A}_{k-1}, \mathbf{A}_{k+1}, \dots, \mathbf{A}_{d+1}, \text{idx}, \text{wgt})$  ▷  $\tilde{\mathbf{Z}} \in \mathbb{R}^{\bar{s} \times r}$ 
10:         $\tilde{\mathbf{X}} \leftarrow \text{TNSRSAMP}(\mathcal{X}, k, \text{idx}, \text{wgt})$  ▷  $\tilde{\mathbf{X}} \in \mathbb{R}^{\bar{s} \times n_k}$ 
11:         $\mathbf{A}_k \leftarrow \arg \min_{\mathbf{B} \in \mathbb{R}^{n_k \times r}} \|\tilde{\mathbf{Z}}\mathbf{B}^\top - \tilde{\mathbf{X}}^\top\|$ 
12:         $\mathbf{p}_k \leftarrow \ell(\mathbf{A}_k)/r$ 
13:      end for
14:    end for
15:    Compute fit (exact or approximate) ▷ Computed only after each epoch
16:    until fit has not improved by more than  $\text{tol}$  for  $\pi$  subsequent epochs
17:    return  $[\mathbf{A}_1, \mathbf{A}_2, \dots, \mathbf{A}_{d+1}]$ 
18: end function

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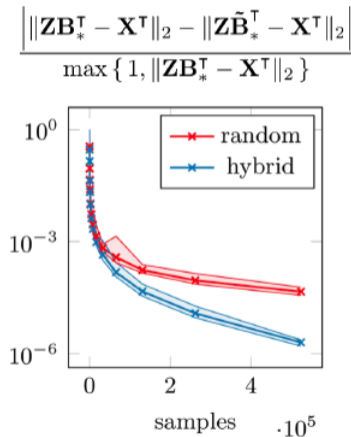
Hybrid leverage score sampling

Single Least Squares Problem with $N = 46M$ rows, $r = 10$ columns, $n = 183$ right-hand sides



$$\tilde{B}_* \equiv \arg \min_{B \in \mathbb{R}^r} \|\Omega Z B^T - \Omega X^T\|_2^2$$

$$B_* \equiv \arg \min_{B \in \mathbb{R}^r} \|Z B^T - X^T\|_2^2$$



Matlab Demo