CSE 392/CS 395T/M 397C: Matrix and Tensor Algorithms for Data

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University of Texas, Austin Spring 2025 Lecture 4: Matrix factorizations I - QR, SVD

Outline

Orthogonality

2 QR Decomposition

3 Singular Value Decomposition

Orthogonality

- Two vectors \boldsymbol{u} and \boldsymbol{v} are orthogonal if $\langle \boldsymbol{u}, \boldsymbol{v} \rangle = 0$.
- A set of vectors $\{\boldsymbol{u}_1, \dots, \boldsymbol{u}_d\}$ is orthogonal if $\langle \boldsymbol{u}_i, \boldsymbol{u}_j \rangle = 0$ for $i \neq j$; and orthonormal if $\langle \boldsymbol{u}_i, \boldsymbol{u}_j \rangle = \delta_{ij}$ for i = j.
- $U \in \mathbb{R}^{n \times d}$ is orthonormal if $U^{\top}U = I$. If U is square, then it is orthogonal (or unitary if complex), and $UU^{\top} = I$.
- Orthonormal matrices preserve norms: $\|Uy\|_2 = \|y\|_2$.

Projectors

Projection matrix: A symmetric matrix P of the form $P = UU^{\top}$ is an orthogonal projection matrix, with:

- $P^2 = P$.
- If P is a (orthogonal) projection matrix, then:

$$ar{m{P}} = m{I} - m{P}$$

is also a projection matrix.

• If U is an orthonormal basis of $\mathbb{X} \subseteq \mathbb{R}^n$, then:

$$Ran(\mathbf{P}) = \mathbb{X}$$
, and $Ran(\mathbf{I} - \mathbf{P}) = Null(\mathbf{P}) = \mathbb{X}^{\perp}$

Question: $P\bar{P} = ?$

Subspaces of a matrix

Let $\mathbf{A} \in \mathbb{R}^{n \times d}$ and consider $Ran(\mathbf{A})^{\perp}$, then:

$$Ran(\mathbf{A})^{\perp} = Null(\mathbf{A}^{\top})$$

Proof: Any $x \in Ran(A)^{\perp}$ iff $\langle Ay, x \rangle = 0$ for all y.

This is same as $\langle \boldsymbol{y}, \boldsymbol{A}^{\top} \boldsymbol{x} \rangle = 0$ for all \boldsymbol{y} .

Similarly, we also have:

$$Ran(\mathbf{A}^{\top}) = Null(\mathbf{A})^{\perp}$$

Thus:

$$\mathbb{R}^{n} = Ran(\mathbf{A}) \oplus Null(\mathbf{A}^{\top})$$

$$\mathbb{R}^{d} = Ran(\mathbf{A}^{\top}) \oplus Null(\mathbf{A})$$

Finding an orthonormal basis of a subspace

- Goal: Find vector in span(A) closest to some vector b.
- Much easier with an orthonormal basis for $span(\mathbf{A})$.

Given $A = [a_1, ..., a_d]$, compute $Q = [q_1, ..., q_d]$ which has orthonormal columns and s.t. span(Q) = span(A).

Each column of A must be a linear combination of certain columns of Q.

Gram-Schmidt process: Compute Q so that a_j (j column of A) is a linear combination of the first j columns of Q.

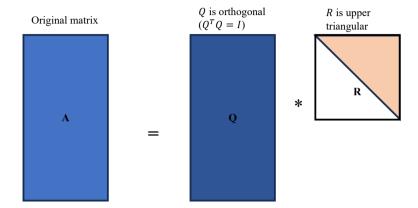
The QR Decomposition

Given $\mathbf{A} \in \mathbb{R}^{n \times d}$ with $n \geq d$, and rank $(\mathbf{A}) = d$, there is a $\mathbf{Q} \in \mathbb{R}^{n \times d}$ and $\mathbf{R} \in \mathbb{R}^{d \times d}$, s.t.

- ullet A=QR
- Q has orthonormal columns, $Q^{\top}Q = I$.
- \mathbf{R} is upper triangular, $r_{ij} = 0$ for i > j.

We have $span(\mathbf{Q}) = span(\mathbf{A})$, the columns of \mathbf{Q} are an orthonormal basis of $span(\mathbf{A})$.

Question: What is the computational cost of QR?



Least squares using QR

• Recall: In the least-squares regression problem, assuming $n \ge d$, we solve:

$$oldsymbol{x}^* = \min_{oldsymbol{x} \in \mathbb{R}^d} \|oldsymbol{A}oldsymbol{x} - oldsymbol{b}\|_2^2.$$

- If A is full rank then we compute A = QR.
- The normal equation can be written as:

$$egin{aligned} oldsymbol{A}^ op oldsymbol{A} oldsymbol{A} & oldsymbol{A}^ op oldsymbol{Q} oldsymbol{R} oldsymbol{A} & oldsymbol{R}^ op oldsymbol{Q} oldsymbol{R} oldsymbol{X} & oldsymbol{A} oldsymbol{C} oldsymbol{A} oldsymbo$$

• Therefore,

$$\boldsymbol{x}^* = \boldsymbol{R}^{-1} \boldsymbol{Q}^\top \boldsymbol{b}.$$

Note that R is non-singular.

- Alternatively, recall that $span(\mathbf{Q}) = span(\mathbf{A})$.
- We know that $\|\mathbf{A}\mathbf{x} \mathbf{b}\|_2$ is minimum when $\mathbf{A}\mathbf{x} \mathbf{b} \perp span(\mathbf{Q})$.
- This implies what?

As a rule it is not a good idea to form $A^{\top}A$ and solve the normal equations. Methods using the QR factorization are better. Why?

QR factorization is also used in direct solvers of linear system Ax = b.

The Singular Value Decomposition

SVD

For any matrix $A \in \mathbb{R}^{n \times d}$ there exist unitary matrices $U \in \mathbb{R}^{n \times n}$ and $V \in \mathbb{R}^{d \times d}$ such that

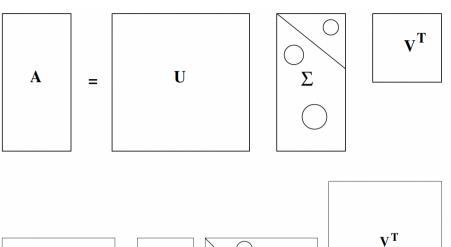
$$\boldsymbol{A} = \boldsymbol{U} \Sigma \boldsymbol{V}^{\top}$$

where Σ is a diagonal matrix with entries $\sigma_i \geq 0$.

$$\sigma_1 \ge \sigma_2 \ge \cdots \ge \sigma_p \text{ with } p = \min(n, d)$$

Let $\sigma_1 = \|\mathbf{A}\|_2 = \max_{\mathbf{x}, \|\mathbf{x}\|=1} \|\mathbf{A}\mathbf{x}\|_2$. There exists a pair of unit vectors such that

$$\mathbf{A}\mathbf{v}_1 = \sigma_1\mathbf{u}_1.$$



A

ι



Thin SVD

• In the first case, suppose, we can write

$$oldsymbol{A} = \left[oldsymbol{U}_1 \, oldsymbol{U}_2
ight] egin{bmatrix} \Sigma \ 0 \end{bmatrix} oldsymbol{V}^ op,$$

where $U_1 \in \mathbb{R}^{n \times d}$ and $U_2 \in \mathbb{R}^{n \times n - d}$. Then,

$$\boldsymbol{A} = \boldsymbol{U}_1 \boldsymbol{\Sigma}_1 \boldsymbol{V}^{\top},$$

where $\Sigma_1, \mathbf{V} \in \mathbb{R}^{d \times d}$.

• Referred to as thin or economical SVD.

Question: How to compute the thin SVD of \boldsymbol{A} from its QR factorization?

SVD Properties

Suppose

$$\sigma_1 \geq \sigma_2 \geq \cdots \geq \sigma_r > 0$$
 and $\sigma_{r+1} = \cdots = \sigma_{\min(n,d)} = 0$

Then:

- $rank(\mathbf{A}) = r = number of nonzero singular values.$
- $Ran(\mathbf{A}) = span\{\mathbf{u}_1, \mathbf{u}_2, \dots, \mathbf{u}_r\}$
- $Null(\mathbf{A}^{\top}) = span\{\mathbf{u}_{r+1}, \mathbf{u}_{r+2}, \dots, \mathbf{u}_n\}$
- $Ran(\mathbf{A}^{\top}) = span\{\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_r\}$
- $Null(\mathbf{A}) = span\{\mathbf{v}_{r+1}, \mathbf{v}_{r+1}, \dots, \mathbf{v}_d\}$

SVD Properties II

 \bullet A matrix \boldsymbol{A} admits the SVD expansion

$$oldsymbol{A} = \sum_{i=1}^r \sigma_i oldsymbol{u}_i oldsymbol{v}_i^ op$$

- $\|A\|_2 = \sigma_1 = \text{largest singular value}$.
- $\bullet \|\mathbf{A}\|_F = \sqrt{\sum_{i=1}^r \sigma_i^2}.$

Eckart-Young-Mirsky Theorem

For any matrix $\boldsymbol{A} \in \mathbb{R}^{n \times d}$ with rank r, let $k \leq r$ and $\boldsymbol{A}_k = \sum_{i=1}^k \sigma_i \boldsymbol{u}_i \boldsymbol{v}_i^{\top}$ then

$$\min_{\boldsymbol{B}: \operatorname{rank}(\boldsymbol{B}) = k} \|\boldsymbol{A} - \boldsymbol{B}\|_2 = \|\boldsymbol{A} - \boldsymbol{A}_k\|_2 = \sigma_{k+1}.$$

Pseudo-inverse

• Given $\mathbf{A} = \mathbf{U} \Sigma \mathbf{V}^{\top}$, we rewrite it as:

$$oldsymbol{A} = \left[oldsymbol{U}_1 \, oldsymbol{U}_2
ight] \left[egin{matrix} \Sigma_1 \, 0 \ 0 & 0 \end{smallmatrix}
ight] \left[oldsymbol{V}_1^ op \ oldsymbol{V}_2^ op \end{smallmatrix}
ight] = oldsymbol{U}_1 \Sigma_1 oldsymbol{V}_1^ op$$

• Then the pseudo inverse of A is:

$$oldsymbol{A}^\dagger = oldsymbol{V}_1 \Sigma_1^{-1} oldsymbol{U}_1^ op = \sum_{i=1}^r rac{1}{\sigma_i} oldsymbol{v}_i oldsymbol{u}_i^ op$$

• The pseudo-inverse of A is the mapping from a vector b to the (unique) Minimum Norm solution of the LS problem: $\min_{x \in \mathbb{R}^d} ||Ax - b||_2^2$.

$$\boldsymbol{x} = (\boldsymbol{A}^{\top} \boldsymbol{A})^{-1} \boldsymbol{A}^{\top} \boldsymbol{b} = \boldsymbol{A}^{\dagger} \boldsymbol{b}.$$

- Let us express solution x in basis V as: $x = Vy = [V_1, V_2] \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$.
- Then left multiply by U^{\top} to get:

$$\|oldsymbol{A}oldsymbol{x} - oldsymbol{b}\|_2^2 = \left\| egin{bmatrix} \Sigma_1 \ 0 \end{bmatrix} egin{bmatrix} oldsymbol{y}_1 \ oldsymbol{y}_2 \end{bmatrix} - egin{bmatrix} oldsymbol{U}_1^ op oldsymbol{b} \ oldsymbol{U}_2^ op oldsymbol{b} \end{bmatrix}
ight\|_2^2$$

- Let us find all possible solutions in terms of $y = [y_1; y_2]$.
- From above, we have $\mathbf{y}_1 = \Sigma_1^{-1} \mathbf{U}_1^{\mathsf{T}} \mathbf{b}$ and \mathbf{y}_2 can be anything.
- Then,

$$egin{array}{lll} oldsymbol{x} &=& [oldsymbol{V}_1, \, oldsymbol{V}_2] egin{array}{lll} oldsymbol{y}_1 &=& oldsymbol{V}_1 \Sigma_1^{-1} oldsymbol{U}_1^ op oldsymbol{b} + oldsymbol{V}_2 oldsymbol{y}_2 \ &=& oldsymbol{A}^\dagger oldsymbol{b} + oldsymbol{V}_2 oldsymbol{y}_2. \end{array}$$

- We know that $\mathbf{A}^{\dagger}\mathbf{b} \in Ran(\mathbf{A}^{\top})$ and $\mathbf{V}_{2}\mathbf{y}_{2} \in Null(\mathbf{A})$.
- Therefore: least-squares solutions are all of the form:

$$A^{\dagger}b + w$$
 where $w \in Null(A)$.

- We obtain the smallest norm when $\mathbf{w} = 0$.
- The Minimum Norm solution of the LS problem: $\min_{x \in \mathbb{R}^d} \|Ax b\|_2^2$ is:

$$\boldsymbol{x}_{LS} = \boldsymbol{V}_1 \Sigma_1^{-1} \boldsymbol{U}_1^{\top} \boldsymbol{b} = \boldsymbol{A}^{\dagger} \boldsymbol{b}.$$

Moore-Penrose Inverse

The pseudo-inverse of $\mathbf{A} \in \mathbb{R}^{n \times d}$ is given by

$$oldsymbol{A}^\dagger = oldsymbol{V} egin{bmatrix} \Sigma_1^{-1} \ 0 \end{bmatrix} oldsymbol{U}^ op = \sum_{i=1}^r rac{1}{\sigma_i} oldsymbol{v}_i oldsymbol{u}_i^ op \end{bmatrix}$$

Properties:

- ullet $oldsymbol{A}oldsymbol{A}^\daggeroldsymbol{A}$
- $\mathbf{A}^{\dagger}\mathbf{A} = \mathbf{I}$ when rank $(\mathbf{A}) = d$, and $\mathbf{A}^{\dagger} = \mathbf{A}^{-1}$ if \mathbf{A} is invertible.
- Left inverse: $A^{\dagger} = (A^{\top}A)^{-1}A^{\top}$, when $n \geq d$, and A is full rank.
- Right inverse: $A^{\dagger} = A^{\top} (AA^{\top})^{-1}$, when $n \leq d$, and A is full rank.

Exercises

- AA^{\dagger} is a projector onto which space?
- $A^{\dagger}A$ is a projector onto which space?
- For orthonormal U, show that $U^{\top}b = \arg\min_{x} \|Ux b\|_2$. (You might use the normal equations, or the Pythagorean theorem as stated for projections.)
- For orthonormal U, show that $P_U b$ is the closest vector to b in span(U).
- Show that $\mathbf{A}^{\dagger} = (\mathbf{A}^{\top} \mathbf{A})^{\dagger} \mathbf{A}^{\top}$.
- Show for symmetric P that $PP = P \implies P = UU^{\top}$ for some orthonormal U.

Questions?